## On Spatial Reasoning with Description Logics

- Motivation
- The family of $\mathcal{A L C} \mathcal{I}_{\mathcal{R C C}}$ logics
- Work in progress
- What we know
- What we don't know
- Future Work
- We want a DL for "qualitative composition-table based spatial reasoning" in the style of $\operatorname{ALC} \mathcal{C P}\left(\mathcal{S}_{2}\right)$, but without syntax-restrictions (if possible)
- With roles corresponding to RCC relationships
- Cohn '93: Multi-modal spatial logic with " $\square_{R}, \diamond_{R}$ " for each RCC-relationship $R$
- Purely relational semantics
(no truly spatial interpretations yet)
- Related to Relation Algebras, but weaker semantics (e.g., our models must not necessarily be representations of finite relation algebras)


## The $\mathcal{A L C I}_{\mathcal{R C C}}$-family

- We are considering this problem in a DL-setting
- In contrast to previous work: inverse roles
- $\mathcal{A L C I}$ with disjoint roles and global role axioms of the form $S \circ T \sqsubseteq R_{1} \sqcup \cdots \sqcup R_{n}$
- Semantics:

$$
\begin{gathered}
\mathcal{I} \models S \circ T \sqsubseteq R_{1} \sqcup \cdots \sqcup R_{n} \text { iff } \\
S^{\mathcal{I}} \circ T^{\mathcal{I}} \subseteq R_{1}^{\mathcal{I}} \cup \cdots \cup R_{n}^{\mathcal{I}}
\end{gathered}
$$

- With role boxes corresponding to RCC1, RCC2, RCC3, RCC5, RCC8: " $\mathcal{A L C} \mathcal{I}_{\mathcal{R C c}}$-family", $\mathcal{A L C I}_{\mathcal{R C C}_{1}}, \mathcal{A L C I}_{\mathcal{R C C}_{2}}, \ldots, \mathcal{A L C I}_{\mathcal{R C C}}$
- With arbitrary role boxes: undecidable (representability of Relation Algebras is undecidable)


## Composition Table Based Reasoning: RCC8 Slide 4


$a \quad b \quad c$
$D C(a, c)$

$E C(a, c)$

$\operatorname{PO}(a, c) \operatorname{TPP}(a, c)$

$\operatorname{TPPI}(a, c)$

Given $\operatorname{EC}(a, b), E C(b, c)$, what do we know about the relationship between $a$ and $c$ ? Lookup $E C \circ E C$ in the RCC8 composition-table:

$$
\begin{aligned}
& \forall x, y, z: E C(x, y) \wedge E C(y, z) \Rightarrow \\
& (D C(x, z) \vee E C(x, z) \vee P O(x, z) \vee \\
& \operatorname{TPP}(x, z) \vee \operatorname{TPPI}(x, z))
\end{aligned}
$$

$E C \circ E C \sqsubseteq D C \sqcup E C \sqcup P O \sqcup T P P \sqcup T P P I$

## Qualitative Spatial Reasoning Example

$$
\begin{array}{lrl}
\text { circle } & \doteq & \text { figure } \\
\text { figure_touching_a_figure } & \doteq & \text { figure } \sqcap \exists \text { EC.figure } \\
\text { special_figure } & \doteq & \text { figure } \sqcap \\
& \forall P O . \neg \text { figure } \sqcap \\
& \forall N T P P I . \neg \text { figure } \sqcap \\
& \forall T P P I . \neg \text { circle } \sqcap \\
& \exists T P P I .(\text { figure } \sqcap \exists E C . \text { circle })
\end{array}
$$

```
special_figure \sqsubseteq figure_touching_a_figure iff
```

figure $\sqcap \forall P O . \neg$ figure $\sqcap \forall N T P P I . \neg$ figure $\sqcap \forall T P P I . \neg$ circle $\sqcap$
$\exists T P P I .($ figure $\sqcap \exists E C . c i r c l e) ~ \sqcap \neg(f i g u r e ~ \sqcap \exists E C . f i g u r e)$ is unsatisfiable w.r.t.
$\mathfrak{R}=\{\ldots, T P P I \circ E C \sqsubseteq E C \sqcup P O \sqcup T P P I \sqcup N T P P I, \ldots\}$
$\boldsymbol{D}$ $\left(\partial^{\prime} \mathrm{D}\right.$ IddILN ${ }^{\prime}\left(\partial^{\prime} q\right)$ DG


## $\mathcal{A L C I}_{\mathcal{R C C} 1}$

- "RCC1": Only one spatial role $S R$, "spatially related"
- Composition table: $\{S R \circ S R \rightarrow S R\}$
- $S R$ is an equivalence relation
- Equivalent to modal logic "S5"
- "S5" reduction principles:
$\diamond p \equiv \square \diamond p, \square p \equiv \diamond \square p, \diamond p \equiv \diamond \diamond p, \square p \equiv \square \square p$
$\Rightarrow$ nested occurrences of modalities can be flattened
- NP-complete satisfiability problem


## $\mathcal{A L C}^{\mathcal{R C C} 2}$

## Slide 8



- "RCC2": reflexive, symmetric role $O=$ "overlap", irreflexive and symmetric role $D R=$ "discrete from"
- Models are fairly trivial: each complete random graph with $\operatorname{Id}\left(\Delta^{\mathcal{I}}\right) \subseteq O^{\mathcal{I}}$ is a model of the role box
- Instead of reduction principles, we have axioms like $\exists O . C \Rightarrow \forall O .(C \sqcup \exists\{O, D R\} . C) \sqcap \forall D R . \exists\{O, D R\} . C)$
- Complexity?


## $\mathcal{A L C}_{\mathcal{R C C} 3} \ldots \mathcal{A L C I}_{\mathcal{R C C} 8}$ : Role Constraints $\quad$ Slide 9

- $\geq \mathcal{A L C} \mathcal{I}_{\mathcal{R C C} 3}:$ There is a special role $E Q$
- Semantics:
- "Weak": $\operatorname{Id}\left(\Delta^{\mathcal{I}}\right) \subseteq E Q^{\mathcal{I}} \Rightarrow$ "Equality" ("EQ" is congruence relation for roles)
- "Strong": $\operatorname{Id}\left(\Delta^{\mathcal{I}}\right)=E Q^{\mathcal{I}} \Rightarrow$ "Identity" (as in Relation Algebras: "EQ" is congruence relation for roles and concepts)
- Further constraints, according to the RCC table
- Reflexiveness, e.g. "Overlap"
- Symmetry, e.g. "Externally Connected"
- Anti-symmetry and irreflexiveness, e.g. "Proper Part"


## $\mathcal{A L C I}_{\mathcal{R C C} 3}$ is Decidable

| $\circ$ | $D R(a, b)$ | $O N E(a, b)$ | $E Q(a, b)$ |
| :---: | :---: | :---: | :---: |
| $D R(b, c)$ | $*$ | $\{D R, O N E\}$ | $D R$ |
| $O N E(b, c)$ | $\{D R, O N E\}$ | $*$ | $O N E$ |
| $E Q(b, c)$ | $D R$ | $O N E$ | $E Q$ |

With the strong EQ semantics, an easy translation into $\mathcal{F}_{2}(=)$ can be given: simply replace "EQ" in $C$ with " $=$ "

$$
\begin{aligned}
\phi_{x}\left(C_{E Q \leftarrow=}\right) \wedge & \forall x, y: \operatorname{DR}(x, y) \oplus \operatorname{ONE}(x, y) \oplus x=y \wedge \\
& \forall x, y: \operatorname{DR}(x, y) \Leftrightarrow \operatorname{DR}(y, x) \wedge \\
& \forall x, y: \operatorname{ONE}(x, y) \Leftrightarrow \operatorname{ONE}(y, x)
\end{aligned}
$$

## $\mathcal{A L C I}_{\mathcal{R C C}_{3}}$ is Decidable (2)

- With the weak $E Q$-semantics, things are not so obvious
- Not every complete, $\{D R, O N E, E Q\}$-edge-colored graph is a model for the role box axioms
- We have to verify that

$$
\begin{aligned}
\forall x, y, z: E Q(x, z) \Leftrightarrow & D R(x, y) \wedge D R(y, z) \oplus \\
& O N E(x, y) \wedge O N E(y, z) \oplus \\
& E Q(x, y) \wedge E Q(y, z)
\end{aligned}
$$

holds, using only two variables

- Idea: use " $=$ " to enforce network consistency, but take care of the fact that "="-connected objects may have different propositional descriptions


## $\mathcal{A L C I}_{\mathcal{R C C}_{3}}$ is Decidable (3)



- Nodes in $E Q$-clique have equivalent modal point of view
- May have different propositional descriptions
- Left structure needs three, right structure only two variables for description


## $\mathcal{A L C I}_{\mathcal{R C C} 5}$ \& $\mathcal{A L C}^{\mathcal{R C C} 8}$ <br> Slide 13

- No finite model property

- $\mathcal{A L C I}_{\mathcal{R C C}_{5}}: P \mathbf{P P}, \mathbf{P P I}$
- $\mathcal{A L C I}_{\text {RCC } 8}$ : $T P P, T P P I, N T P P, N T P P I$
- $\mathcal{A L C I}_{\mathcal{R C C} 8}$ somehow allows the distinction of a role and its transitive orbit ( $\rightarrow$ "PDL binary counter" concept possible)
- This seems to be impossible in $\mathcal{A L C I}_{\mathcal{R C C} 5}$


## The Concept even_odd_chain

```
\(e v e n \_o d d \_c h a i n=d e f\)
    even \(\sqcap\)
    ( \(\exists\) TPPI. \(\exists\) TPPI.丁) П
    \((\) even \(\Rightarrow \forall\) TPPI.odd) \(\sqcap\)
    (odd \(\Rightarrow \forall\) TPPI.even) \(\sqcap\)
    ( \(\forall\) NTPPI.( (even \(\Rightarrow \forall\) TPPI.odd) \(\sqcap\)
        \((\) odd \(\Rightarrow \forall\) TPPI.even \())\) ) \(\sqcap\)
    ( \(\forall\) TPPI. \((\) (even \(\Rightarrow \forall\) TPPI.odd \() \sqcap\)
        \((\) odd \(\Rightarrow \forall\) TPPI.even \())\) ) \(\sqcap\)
    ( \(\forall N T P P I . \exists T P P I . \top)\)
\(\left(\left(T P P I^{\mathcal{I}}\right)^{+}-\boldsymbol{T P P} I^{\mathcal{I}}\right) \subseteq N T P P I^{\mathcal{I}}\)
```


## Is it Possible to Represent Grids?



Michael Wessel, April 2002

## Is it Possible to Represent Grids? (2)

Slide 16


Michael Wessel, April 2002

## Is it Possible to Represent Grids? (3)



Even though infinite grid-like models exists, we found no way to enforce the coincidence of the
$\boldsymbol{x} \circ \boldsymbol{y}$ - and $\boldsymbol{y} \circ \boldsymbol{x}-$
successors.

## Finite Model Reasoning with $\mathcal{A L C I}_{\mathcal{R C C}_{5}}$ ?

- $\mathcal{A L C} \mathcal{I}_{\text {RCC } 5}$ contains the "proper part" role $P P$
- Question:

Suppose we disallow the use of $P P$ in concepts then, do we have the finite model property back?

- Answer: No! Counter example:

```
\existsDR.丁 п
\forallR.( \existsPO.\existsDR.C \sqcap
    PO.\negC \sqcap
    DR.\negC)
```

$\Rightarrow$ There does not seem to be a way to tell, syntactically, whether a concept admits a finite model

## Future Work

- Check out results from "Algebraic Logic"
- Representability of Relation Algebras (RAs) is, generally, undecidable
* There can not be a (decidable) $\mathcal{A} \mathcal{L C} \mathcal{I}_{\mathcal{R A}}$ with arbitrary role boxes
- So is the equational theory of arbitrary RAs
- Decidable classes of (relation) algebras that are useful for spatial reasoning with DLs?
- Multi-dimensional modal logics
- Arrow-logic

