

On Terminological Default Reasoning about Spatial Information

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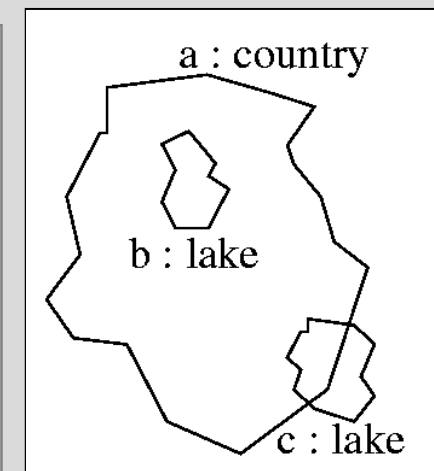
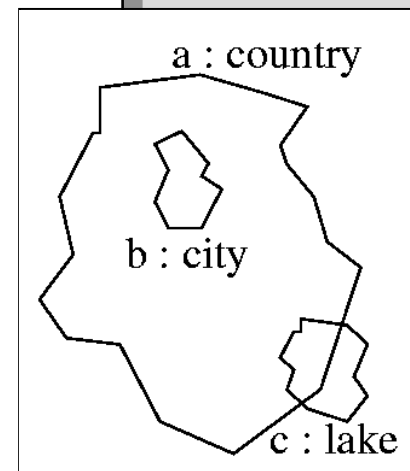
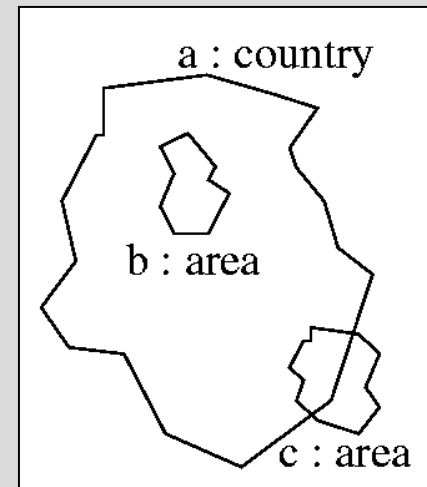
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- ❑ Motivation
- ❑ Preliminaries
 - ⇒ spatioterminological reasoning with $ALCRP(S_2)$ [Lutz, Haarslev & Möller]
 - ⇒ nonmonotonic reasoning with Reiter's default logic
 - terminological default rules / theories [Baader & Hollunder]
- ❑ Spatioterminological default reasoning
 - ⇒ on computing extensions
- ❑ Conclusion & future work

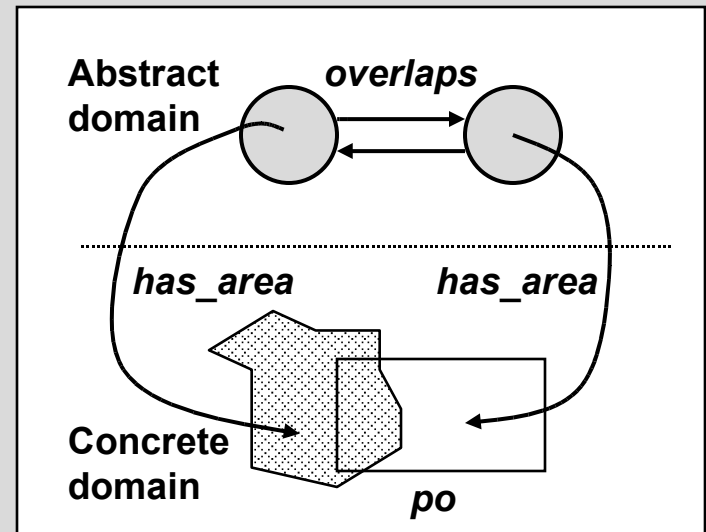
Motivation: Incomplete Spatioterminological Knowledge

- ❑ Combination of terminological, spatial & default reasoning techniques
 - Geographic Information Systems (GIS)
- ❑ Terminological knowledge
 - **capital_city** , **city**
- ❑ Spatial knowledge
 - properties of spatial relationships, eg. **tp** (contains) is **transitive**
- ❑ Spatioterminological knowledge
 - a **city** is **contained** within exactly one **country**
 - two **countries** never **overlap** each other
- ❑ Default knowledge
 - data augmentation / completion
 - „b“ could possibly be a **city** or a **lake**, but not both (disjoint concepts)
 - ABox realization would not work



Spatioterminological Reasoning with $ALCRP(S_2)$

- $ALCRP(\mathcal{D})$ extends $ALC(\mathcal{D})$
 - ⇒ $ALCRP(\mathcal{D}) = ALC(\mathcal{D}) +$ role-forming predicate-based operator
 - ⇒ decidable for restricted concept terms
 - ⇒ restrictedness closed under negation
- $ALCRP(S_2)$
 - ⇒ admissible concrete domain S_2 , regular closed subsets of \mathbb{R}^2 , called **regions**, with **RCC8 predicates**
 - ⇒ properties of relationships captured by concrete domain, e.g. transitivity of **tpp**
 - ⇒ RCC8 predicates
 - **dc, ec, po, tpp, ntp, tppi, ntpi, eq**
 - ⇒ defined roles, TBox axioms
 - $inside \doteq \exists(has_area)(has_area).tpp-ntp$
 - $contains \doteq \exists(has_area)(has_area).tppi-ntppi$
 - $overlaps \doteq \exists(has_area)(has_area).po$



$$(\exists(u_1, \dots, u_n)(v_1, \dots, v_m).P)^I := \{(a, b) \in \Delta_I \times \Delta_I \mid \exists x_1, \dots, x_n, y_1, \dots, y_m \in \Delta_D : (a, x_1) \in u_1^I, \dots, (a, x_n) \in u_n^I, (b, y_1) \in v_1^I, \dots, (b, y_m) \in v_m^I, (x_1, \dots, x_n, y_1, \dots, y_m) \in P^D\}$$

Spatioterminological Background Knowledge (TBox)

$$\begin{aligned} \text{area} &\doteq \exists \text{has_area.is-region} \\ \text{country_region} &\dot{\sqsubseteq} \neg \text{natural_region} \sqcap \\ &\quad \text{large_scale} \sqcap \text{area} \\ \text{city_region} &\dot{\sqsubseteq} \neg \text{natural_region} \sqcap \\ &\quad \neg \text{large_scale} \sqcap \text{area} \\ \text{lake_region} &\dot{\sqsubseteq} \text{natural_region} \sqcap \text{area} \\ \text{country} &\doteq \text{country_region} \sqcap \\ &\quad \forall \text{contains.} \neg \text{country_region} \sqcap \\ &\quad \forall \text{overlaps.} \neg \text{country_region} \sqcap \\ &\quad \forall \text{inside.} \neg \text{country_region} \\ \text{city} &\doteq \text{city_region} \sqcap \\ &\quad \exists \text{inside.country_region} \\ \text{lake} &\dot{\sqsubseteq} \text{lake_region} \end{aligned}$$

Default Theories & Terminological Default Theories

- ❑ Default rules [Reiter, 1980]
 - ⇒ α prerequisite, β_i justifications, γ conclusion, FOPL formulae
- ❑ Default theory (W, D)
 - ⇒ W = world description
 - ⇒ D = set of defaults
- ❑ Different sets of **extensions** of (W, D)
 - ⇒ **sceptical** vs. **credulous** consequence
- ❑ **Terminological** default theories [Baader & Hollunder, 1991]
 - ⇒ α, β_i, γ concept terms
 - ⇒ W = ABox, D = set of closed default rules
 - ⇒ restricted semantics, no skolemization
 - ⇒ consequence problem decidable
- ❑ Closing concept terms over ABox W
 - ⇒ concept terms \Rightarrow ABox concept membership assertions

$$\frac{\alpha : \beta_1, \beta_2, \dots, \beta_n}{\gamma}$$

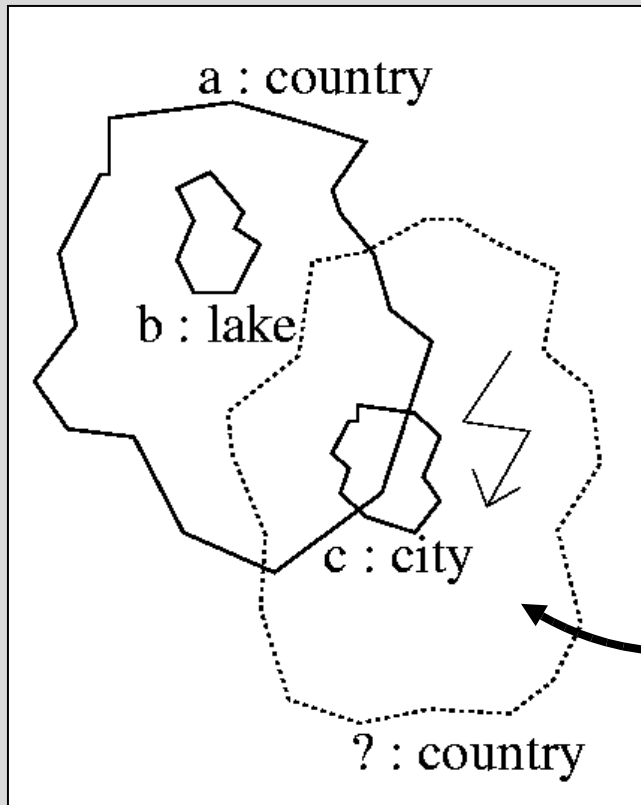
Example

$$\frac{\textit{area} : \textit{country}}{\textit{country}}$$

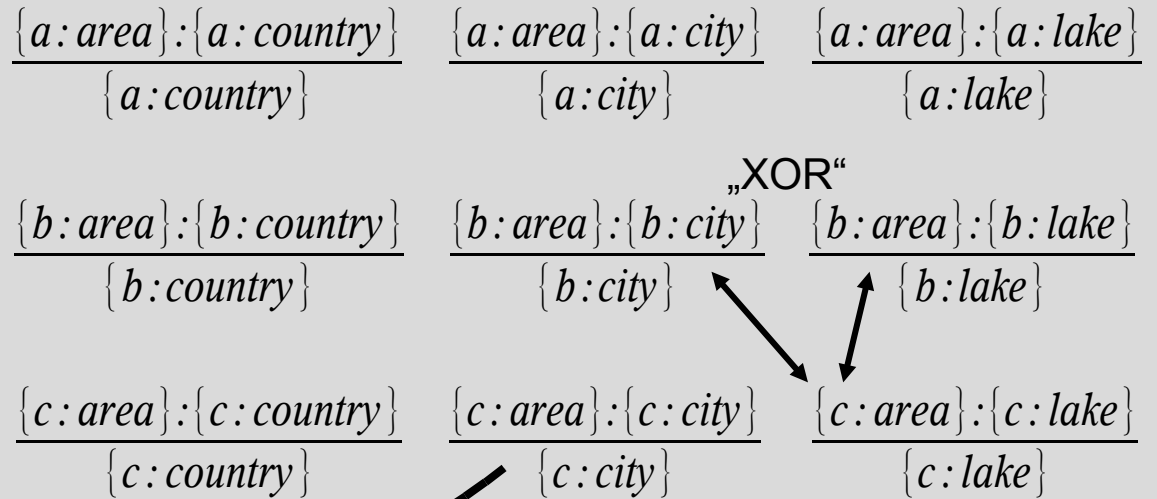
$$\frac{\textit{area} : \textit{city}}{\textit{city}}$$

$$\frac{\textit{area} : \textit{lake}}{\textit{lake}}$$

Closed Defaults



- World description $W =$
 $\{ a : \text{country}, b : \text{area}, c : \text{area},$
 $(a, b) : \text{contains}, (b, a) : \text{inside},$
 $(a, c) : \text{overlaps}, (c, a) : \text{overlaps} \}$



- Closing over W yields 9 closed defaults
- Two extensions
 - $E1 = W \cup \{ \underline{b : \text{city}}, c : \text{lake} \}$
 - $E2 = W \cup \{ \underline{b : \text{lake}}, c : \text{lake} \}$
- $W \cup \{ b : \text{lake}, c : \text{city} \}$
inconsistent, see picture
- 2 sets of „generating defaults“

Spatioterminological Default Theories with „ABox Patterns“

- ❑ We also want to conclude complex role assertions
 - ⇒ $W = \{ \textit{linköping} : \textit{swedish_city}, \textit{sweden} : \textit{country} \}$
 - ⇒ $E = W \cup \{ (\textit{sweden}, \textit{linköping}) : \textit{contains} \}$
 - ⇒ cannot be expressed with concept terms as α, β_i, γ
- ❑ „ABox patterns“
 - ⇒ ABoxes with **variables**, e.g. X, Y, Z
 - ⇒ to be closed over W
 - ⇒ can also refer to specific ABox individuals
- ❑ „Duality“
 - ⇒ use known concept memberships to conclude spatial relationships
 - ⇒ use spatial relationships to deduce concept memberships

$$\frac{\{ X : \textit{swedish}_{city}, \textit{sweden} : \textit{country} \} : \{ (\textit{sweden}, X) : \textit{contains} \}}{\{ (\textit{sweden}, X) : \textit{contains} \}}$$

$$\frac{\{ \textit{linköping} : \textit{swedish}_{city}, \textit{sweden} : \textit{country} \} : \{ (\textit{sweden}, \textit{linköping}) : \textit{contains} \}}{\{ (\textit{sweden}, \textit{linköping}) : \textit{contains} \}}$$

Close 

On Computing Extensions

Let E be a set of closed formulae and (A, D) be a closed default theory. We define $E_0 := A$ and for all $i \geq 0$

$$E_{i+1} := E_i \cup \{ \gamma \mid \alpha : \beta_1, \dots, \beta_n / \gamma \in D, \\ \alpha \in Th(E_i), \\ \neg \beta_1, \dots, \neg \beta_n \notin Th(E) \}.$$

Then, $Th(E)$ is an extension of (A, D) iff

$$Th(E) = \bigcup_{i=0}^{\infty} Th(E_i)$$

$$E_i \quad \alpha, \alpha = \{ a_1, a_2, \dots, a_n \}$$

$$\forall a_i \in \alpha : E_i \quad a_i$$

$$E \quad \neg \beta, \beta = \{ b_1, b_2, \dots, b_n \}$$

$$\forall b_i \in \beta : E \quad \neg b_i$$

- Non-constructive definition, since **$Th(E)$** is already used in each „iteration step“

⇒ however, each extension has the form

$$Th(W \cup Con(D'))$$

for a set of so-called **generating defaults $D', D'' \subseteq D$**

- ⇒ simple „generate & test“ algorithm:
 - „generator“: compute powerset of **$Con(D)$** and „test“ each subset
 - „tester“: use definition to check if candidate is indeed an extension
- ⇒ more efficient algorithms see Baader & Hollunder

- α, β_i, γ are ABoxes
 - $\alpha \in Th(E_i) \Leftrightarrow E_i \quad \alpha$
 - $\neg \beta \notin Th(E) \Leftrightarrow E \quad \neg \beta$

ABox Axiom Entailment

A restricted $\mathcal{ALCRP}(\mathcal{S}_2)$ ABox axiom x
is logically entailed by a restricted $\mathcal{ALCRP}(\mathcal{S}_2)$
ABox A ,

$$A \models x, \quad \text{iff} \quad \begin{cases} x = a : C \longrightarrow \neg SAT(A \cup \{a : \neg C\}) \\ x = (a, b) : \exists(u)(v).P \longrightarrow \\ \neg SAT(A \cup \{(a, b) : \exists(u)(v).\bar{P}\}) \wedge \\ \neg SAT(A \cup \{a : \forall u.\top\}) \wedge \\ \neg SAT(A \cup \{b : \forall v.\top\}) \end{cases}$$

$SAT(A)$ decides the ABox consistency problem for an ABox A , and $u = v = has_area$.

- ABox axiom entailment reduced to ABox consistency (negation necessary)
 - ⇒ α, β_i, γ may only contain
 - concept membership axioms: „instance checking“ problem
 - complex role assertions (cannot be negated, but entailment can be decided)
 - other kinds of axioms possible?

Conclusion & Future Work

❑ Extension to Baader & Hollunder

⇒ ABox patterns

- refer to specific individuals
- complex role assertions

❑ Other kinds of ABox axioms?

- ⇒ however, concept membership assertions and complex role assertions sufficient in our application domain

❑ Default theories with specificity

- ⇒ if more than one default applicable, apply most specific first
- ⇒ additional partial ordering on defaults
- ⇒ S-Extensions instead of R-Extensions

$$d_1 \prec d_2 \Leftrightarrow \alpha(d_1) \quad \alpha(d_2) \wedge \alpha(d_2) \quad \alpha(d_1)$$

❑ Autoepistemic description logics (operators **A** and **K**)?

❑ Implementation

- ⇒ more efficient algorithms for computing extensions